### POW: System-wide Dynamic Reallocation of Limited Power in HPC

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# Power and Energy

- Different but related ideas
  - Rate vs Quantity
- Conversion:

1 Watt = 
$$1 \frac{Joule}{Second}$$

- 1 kWh = 3.6 megajoules
- Infrastructure required for 900 kWh over 1 hour is not the same as 900 kWh over 720 hours.

## HPC System



## Power Scheduler Invariant

#### $\forall t, \sum c_i^t \leq \sum a_i^t \leq L$

- System-wide power limit L
- Number of sockets  $\mathcal{N}$
- A time t
  - Power consumed by socket i at time t
- $c_i^t a_i^t$ Power allocated to socket i at time t

# Naive Static Strategy

#### $\forall t, \sum c_i^t \leq \sum a_i^t \leq L$

- LSystem-wide power limit
- Number of sockets n
- A time t
  - Power consumed by socket i at time t
- $c_i^t a_i^t$ Power allocated to socket i at time t

$$a_i^t = \frac{L}{n} \implies \sum c_i^t \le L$$

#### Power and Runtime



# Job Static Strategy

 $\forall t, \sum c_i^t \leq \sum a_i^t \leq L$ 

- $L \mid$  System-wide power limit
- $n \mid$  Number of sockets
  - A time
- t  $j_+$   $j_n$   $c_i^t$   $a_i^t$
- Maximum power consumed by a socket for job jNumber of sockets in job j
- Power consumed by socket i at time t
- Power allocated to socket i at time t

$$\forall j, \sum j_+ j_n \leq L \implies \sum c_i^t \leq L$$

# Power and Energy



# Naive Dynamic Strategy

 $\forall t, \sum c_i^t \leq \sum a_i^t \leq L$ 

- - $t \mid A time$
- $\begin{array}{c|c} w_i^t & \text{Waste power for socket } i \text{ at time } t \\ c_i^t & \text{Power consumed by socket } i \text{ at time } t \\ a_i^t & \text{Power allocated to socket } i \text{ at time } t \end{array}$

$$\begin{aligned} \forall t, L &= \sum a_i^t \qquad c_i^t + w_i^t = a_i^t \qquad c_i^t \approx c_i^{t+1} \\ w_i^{t+1} &\approx \frac{1}{n} \sum a_i^t - c_i^t \implies a_i^{t+1} \approx c_i^t + w_i^{t+1} \end{aligned}$$

### POWsched

procedure MAIN	
while True do	
GETREADINGS	$\triangleright$ Phase 1
ALLOCDOWN	$\triangleright$ Phase 2
ALLOCUP	$\triangleright$ Phase 3
sleep rest of interval	
end while	
end procedure	

## 50W Static and Dynamic

![](_page_10_Figure_1.jpeg)

![](_page_10_Figure_2.jpeg)

Time

# Static vs Dynamic

Experiment	Runtime	Stddev	Improvement
115W static	278.26	9.57	
115W dynamic	276.24	4.84	0.7%
90W static	284.63	3.20	
90W dynamic	277.13	5.04	2.6%
70W static	323.83	4.90	
70W dynamic	278.02	4.97	14.1%
50W static	407.21	18.00	
50W dynamic	371.92	13.23	8.7%

# In Summary

- Power Optimization != Power Bound Enforcement
- Static power allocation may not be optimal
- Dynamic power reallocation can reduce time to solution

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